

Improved bounds for the zeros of the chromatic polynomial

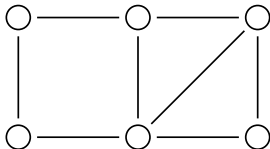
Viresh Patel

Queen Mary, University of London

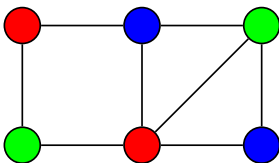
Combinatorial Algorithmic and Probabilistic Aspects of
Partition Functions
CWI Amsterdam

Joint work with
Matthew Jenssen (KCL) and Guus Regts (Amsterdam)

Graph colouring



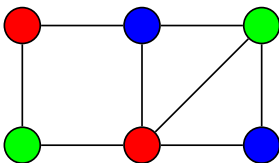
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***k*-colouring of G**

$f : V(G) \rightarrow \{1, \dots, k\}$ s.t. $f(u) \neq f(v)$ whenever $uv \in E(G)$

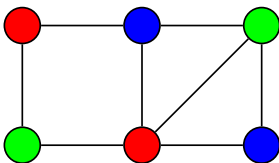
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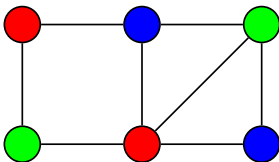
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Chromatic polynomial

$C_G(k) = \#$ k -colourings of G

(Birkhoff 1912)

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- Originally introduced to approach four colour problem
- Examples
 - $C_{K_r}(k) = k(k-1)(k-2) \cdots (k-r+1)$
 - $C_{n\text{-vertex tree}}(k) = k(k-1)^{n-1}$

Computational counting

Want FPTAS / FPRAS for

$$C_G(k) \text{ for } k \in \mathbb{N}$$

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- exactly $\forall z \in \mathbb{C} \setminus \{0, 1, 2\}$. Jaeger-Vertigan-Welsh 1990
- approximately $\forall z$ s.t. $|z - 1| > 1$ Bencs-Huijben-Regts '22

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$C_G(k)$ is NP-hard to approximate for “most” $k \leq \Delta(G)$

[GSV15]

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Conjecture \exists FPTAS for $C_G(k)$ provided $k > \Delta(G)$

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For $k \in \mathbb{N}$

- FPRAS for $k \geq 2\Delta$ Jerrum 1994
- FPRAS for $k > (\frac{11}{6} - \varepsilon)\Delta$ Vigoda 2006, CDMPP 2019
- FPTAS for $k \geq 2\Delta$ Liu-Sinclair-Srivastava 2019

Question: (Brenti, Royle, Wagner 1994)

Is there a function $f(k)$ such that $C_G(z) \neq 0$ whenever $\Delta(G) \leq k$ and $|z| \geq f(k)$?

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Improved bounds for small Δ (Fialho-Juliano-Procacci 2024+)

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Conjecture: $C_G(z) \neq 0$ if $\Re(z) > \Delta(G)$ Sokal 2003

Implies

Conjecture: \exists FPTAS for $C_G(k)$ provided $k > \Delta(G)$
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(sketch)

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(sketch)

Let $G = (V, E)$ with $|V| = n$

$$C_G(z) = a_0 z^n - a_1 z^{n-1} + a_2 z^{n-2} - \dots + (-1)^n a_n$$

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$a_i =$ number of *broken-circuit free sets* of size i in G

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Broken circuit free sets (BCF sets)

Fix an ordering of E . We say $A \subseteq E$ is broken-circuit free if

- A is a forest, and
- each $e \in E \setminus A$ is not the largest edge in the unique cycle of $A + e$ (when it exists)

Note: number of BCF sets is independent of edge ordering!

Example

We work with a simple transformation:

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$$B_G(z) = a_0 + a_1 z + \cdots + a_n z^n = z^n C_G(-z^{-1}) = \sum_{F \subseteq E \text{ BCF}} z^{|F|}$$

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Enough to show that whenever $\Delta(G) \leq \Delta$ and $|z| \leq 1/K\Delta$

$$\left| \frac{B_G(z)}{B_{G-u}(z)} \right| \in [1 - a, 1 + a]$$

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Enough to complete the induction ...

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For final inequality we introduce and bound

$$T_{G,v_1,v_2}(x) = \sum_{\substack{T \text{ tree:} \\ v_1, v_2 \in V(T)}} x^{|T|} \quad \text{and bound} \quad T_{G,v_1,v_2} \left(\frac{\ln \alpha}{\alpha \Delta} \right) \leq \frac{\alpha \ln \alpha}{\Delta}$$

Conclusion

$C_G(z) \neq 0$ whenever $\Delta(G) \leq \Delta$ and $|z| \geq 5.93\Delta$

- Try to improve on 5.93
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$K_3 = 5.93$	$K_4 = 5.23$	$K_5 = 4.87$
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If $\text{girth}(G) \geq g$ and $\Delta(G) \leq \Delta$, can replace 5.93 with $K_{g,\Delta}$

$K_{3,3} = 4.55$	$K_{3,4} = 4.90$	$K_{3,5} = 5.11$
$K_{4,3} = 3.84$	$K_{5,3} = 3.48$	$K_{100,3} = 2.49$

Further Results

Forest generating polynomial of $G = (V, E)$

$$F_G(z) = \sum_{F \subseteq E \text{ forest}} z^{|F|}$$

- Also called partition function of arboreal gas model
- Our methods extend here, but can go further using a different recursion

$F_G(z) \neq 0$ whenever $\Delta(G) \leq \Delta$ and $|z| \leq 1/(2\Delta)$

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- Cannot replace $1/(2\Delta)$ with $1/\Delta$ due to Δ -multi edge.
- Can we get close to $1/\Delta$?