

Occupancy fractions: graph theory and algorithms

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Workshop on the **Combinatorial**, Algorithmic and Probabilistic aspects of Partition Functions

1. What is the smallest independence number of an n -vertex graph of average degree d ?

A: Turán 1941: $\frac{n}{d+1}$ achieved by a disjoint union of K_{d+1} s

2. Which d -regular n -vertex graph has the most independent sets?

A: Kahn 2001, Zhao 2010: a disjoint union of $K_{d,d}$ s

3. Which graph with given degree sequence has the fewest independent sets?

A: SSSZ 2019: a disjoint union of complete graphs

Workshop on the Combinatorial, Algorithmic and Probabilistic aspects of Partition Functions

1. In expectation, how big is the independent set returned by the random greedy algorithm?

A: Folklore (Alon–Spencer 1992?): at least $\sum_v \frac{1}{d_v+1}$

2. In which graphs can you efficiently approximate the number of independent sets of density $1/8$?

A: D., Perkins 2021: graphs of maximum degree 5

3. When can you efficiently approximate the antiferromagnetic Ising partition function at a given magnetization?

A: Not known, partial knowledge for cubic graphs (D., Leblanc 2025)

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Let $E_G(1)$ be the expected density of a uniform random independent set in G

1. Which graph of maximum degree Δ minimizes $E_G(1)$?

A: Cutler, Radcliffe 2014: $K_{\Delta+1}$

2. Which $(2d-1)$ -regular line graph maximizes $E_G(1)$?

A: DJPR 2017a: $L(K_{d,d})$

3. Over triangle-free graphs of maximum degree Δ what is the infimum of $E_G(1)$?

A: Shearer¹, DJPR 2017b: at least $(1 - o(1)) \frac{\log \Delta}{\Delta}$

¹Unpublished, 1992 SIAM talk

Workshop on the Combinatorial, Algorithmic and Probabilistic aspects of **Partition Functions**

- Let $Z_G(\lambda)$ be the **independence polynomial**:

$$Z_G(\lambda) = \sum_{I \subseteq V(G) \text{ indep.}} \lambda^{|I|}$$

- Let $F_G(\lambda)$ be the **free energy density**:

$$F_G(\lambda) = \frac{1}{|V(G)|} \log Z_G(\lambda)$$

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- Let $\mu_{G,\lambda}(I)$ be the **hard-core model** on G

$$\mu_{G,\lambda}(I) = \frac{\lambda^{|I|}}{Z_G(\lambda)}$$

- Let $E_G(\lambda)$ be the **occupancy fraction** of G :

$$E_G(\lambda) = \frac{1}{|V(G)|} \mathbb{E}_{\mu_{G,\lambda}} |I| = \lambda \frac{\partial}{\partial \lambda} F_G(\lambda) = \frac{1}{|V(G)|} \frac{\lambda Z'_G(\lambda)}{Z_G(\lambda)}$$

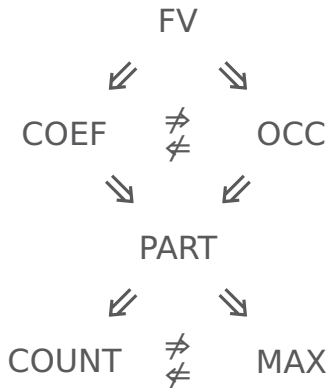
In what ways can partition functions be bigger than others?

$$P(x) = 1 + a_1x + \dots + a_kx^k; Q(x) = 1 + b_1x + \dots + b_kx^k; a_i, b_j \geq 0.$$

Definition

1. Say $P \geq_{\text{MAX}} Q$ if $a_k \geq b_k$
2. Say $P \geq_{\text{COUNT}} Q$ if $\sum_{i=1}^k a_i \geq \sum_{i=1}^k b_i$ $F_P(1) \geq F_Q(1)$
3. Say $P \geq_{\text{PART}} Q$ if $P(x) \geq Q(x)$ for all $x \geq 0$ $F_P(x) \geq F_Q(x)$
4. Say $P \geq_{\text{OCC}} Q$ if $\frac{xP'(x)}{P(x)} \geq \frac{xQ'(x)}{Q(x)}$ for all $x \geq 0$ $E_P(x) \geq E_Q(x)$
5. Say $P \geq_{\text{COEF}} Q$ if $a_i \geq b_i$ for all $1 \leq i \leq k$
6. Say $P \geq_{\text{FV}} Q$ if $b_i a_{i+1} \geq a_i b_{i+1}$ for all $0 \leq i \leq k-1$

Theorem (DJPR 2018)



Lens through which we can see combinatorial results

- Over d -regular graphs K_{d+1} , minimizes $Z_G(\lambda)$ in the FV sense [Cutler, Radcliffe 2014]
- For a given degree sequence, a disjoint union of complete bipartite graphs maximizes $Z_{L(G)}(\lambda)$ in the MAX sense [Bregman 1973]
- For a given degree sequence, a disjoint union of complete bipartite graphs maximizes $Z_G(\lambda)$ in the PART sense [SSSZ 2019]
Fact: strengthening to OCC is *false*
- For a given degree sequence, a disjoint union of complete graphs minimizes $Z_G(\lambda)$ in the PART sense [SSSZ 2019]
Conj. [D., Kang]: strengthening to OCC is *true*

Motivation

Theorem (DJPR 2017)

For a triangle-free graph G of maximum degree Δ ,

$$E_G(\underbrace{1/\log \Delta}_{\lambda}) \geq (1 - o(1)) \frac{\log \Delta}{\Delta}$$

- This is tight as stated by the random Δ -regular graph
- But what about $\lambda \rightarrow \infty$? Then we seem to be a factor 2 off
- We know $E_G(\lambda)$ is increasing in λ , but how fast?
- Perhaps we should study the derivative...

Variance fraction

- Let $V_G(\lambda)$ be the **variance fraction** of G :

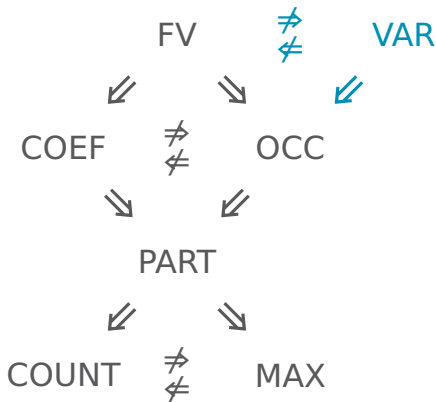
$$\begin{aligned} V_G(\lambda) &= \frac{1}{|V(G)|} \mathbb{V}_{\mu_{G,\lambda}} |I| \\ &= \lambda \frac{\partial}{\partial \lambda} E_G(\lambda) = \frac{1}{|V(G)|} \left(\frac{\lambda^2 Z_G''(\lambda) + \lambda Z_G'(\lambda)}{Z_G(\lambda)} - \frac{\lambda^2 Z_G'(\lambda)^2}{Z_G(\lambda)^2} \right) \end{aligned}$$

- For polynomials P, Q as before, we say that $P \geq_{\text{VAR}} Q$ if

$$\frac{x^2 P''(x) + x P'(x)}{P(x)} - \frac{x^2 P'(x)^2}{P(x)^2} \geq \frac{x^2 Q''(x) + x Q'(x)}{Q(x)} - \frac{x^2 Q'(x)^2}{Q(x)^2}$$

for all $x \geq 0$

General results



Perhaps this is a little surprising, FV seemed pretty strong...

New results for the hard-core model

Theorem (D., Sandhu, Tan 2025)

- Let G be an n -vertex graph. Then for any $0 < \lambda < 1/(2n - 1)$ we have

$$\frac{\lambda}{(1 + n\lambda)^2} = V_{K_n}(\lambda) \leq V_G(\lambda) \leq V_{\overline{K_n}}(\lambda) = \frac{\lambda}{(1 + \lambda)^2},$$

and the upper bound holds up to $\lambda \leq 1/n$.

(not quite VAR due to upper bound on λ)

- It is true that

$$Z_{3K_2+K_3}(\lambda) \leq_{\text{VAR}} Z_{3K_{1,2}}(\lambda)$$

(first nontrivial case of deg-seq VAR lower bound)

New results for the hard-core model

Theorem (D., Sandhu, Tan 2025)

Let G be an n -vertex graph with maximum degree Δ . Then for any $0 < \lambda < 3/(\Delta + 1)^2$ we have

$$\frac{1}{n} \sum_{u \in V(G)} \frac{\lambda}{1 + (d_u + 1)\lambda} = \frac{1}{n} \sum_{u \in V(G)} E_{K_{d_u+1}}(\lambda) \leq E_G(\lambda).$$

(progress on OCC strengthening of SSSZ deg-seq PART bound)

Note that the bound holds in the limit $\lambda \rightarrow \infty$ because it's equivalent to the [Caro–Wei theorem](#) (1978, 1981)

Claim: the proof techniques are interesting

- For VAR: formulate a **higher-order** version of occupancy method techniques which yielded many results for OCC
- Objective is now quadratic instead of linear
- We showed that it **can** work, but it's unclear how sharp these local methods are. Is variance well-explained locally?
- For deg-seq OCC: use an induction approach (cf. SSSZ) with **extension of** local occupancy (cf. DJPR, DJKP, DKPS, DK)
- Lots of things are **tricky**, Taylor expansions get out of hand. . .
- $\lambda = O(1/\Delta)$ seems like a serious barrier for our approach
- So far, methods are combined but not 'properly' interwoven so perhaps more can be done

Open problems

- Understand variance for large λ in any setting
- Extend the range of λ in our results
- Entertain an interest in more parameters:

Theorem (Campos–Samotij 2024)

For all graphs G and $\lambda > 0$, $E_G(\lambda) \leq F_G(\lambda) \frac{\lambda}{(1+\lambda) \log(1+\lambda)}$

Stronger than $\overline{K_n}$ maximizing in the OCC sense!

- Take the derivative $\lambda \frac{\partial}{\partial \lambda}$:

Conjecture (D., Sandhu, Tan 2025)

For all graphs G and $\lambda > 0$, $V_G(\lambda) \leq E_G(\lambda) \frac{1}{1+\lambda}$

Stronger than $\overline{K_n}$ maximizing in the VAR sense!

Thank you



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